

# Specification of all Maximal Subsets of the Hamming Space with Respect to Given Diameter \*

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## Abstract

We present a specification of all maximum subsets of the  $n$ -cube  $B^n$  with respect to a given diameter and show relations between this problem and the discrete isoperimetric problem. As a corollary, for any number  $m$ ,  $1 \leq m \leq 2^n$ , we specify an  $m$ -element subset of  $B^n$  with minimal possible diameter. We also present a simple proof for a theorem of Katona on intersecting families [4].

## 1 Introduction

Denote by  $B^n$  the  $n$ -dimensional unit cube with Hamming metric  $\rho(\alpha, \beta)$ ,  $\alpha, \beta \in B^n$ . For a fixed integer  $d \in [0, n]$  consider the problem of constructing a maximum size subset  $A \subseteq B^n$  such that  $\rho(\alpha, \beta) \leq d$  for any  $\alpha, \beta \in A$ . We call such a subset  $A$   $d$ -maximal. The number  $d$  is called the diameter of  $A$  and denoted by  $D(A)$ . Note that in continuous case the ball of radius  $d/2$  has maximum area among all plane figures of diameter  $d$ . One of the  $d$ -maximal subsets of  $B^n$  is found in [1,2]. The main goal of this paper is to specify all  $d$ -maximal subsets of  $B^n$ . Throughout the paper we assume  $d < n$ . Denote by  $S_r^n(\alpha)$  the ball of radius  $r$  centered in  $\alpha \in B^n$ .

### Theorem 1

- (i) If  $d = 2t$  then  $S_t^n(\alpha)$  is a  $d$ -maximal set for any  $\alpha \in B^n$ ;
- (ii) If  $d = 2t + 1$  then  $S_t^n(\alpha) \cup S_t^n(\beta)$  is a  $d$ -maximal set for any  $\alpha, \beta \in B^n$  such that  $\rho(\alpha, \beta) = 1$ .

For  $\alpha = (\alpha_1, \dots, \alpha_n) \in B^n$  denote  $\bar{\alpha} = (\bar{\alpha}_1, \dots, \bar{\alpha}_n)$ , the complementary vector of  $\alpha$ . Partition  $B^n$  into  $2^{n-1}$  pairs  $(\alpha, \bar{\alpha})$  of complementary vectors and denote by  $\mathcal{M}^n$  the collection of all  $2^{n-1}$ -element subsets of  $B^n$ , obtained by choosing exactly one vector from each pair  $(\alpha, \bar{\alpha})$ .

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## Theorem 2

- (i) If  $d < n - 1$  then there are no  $d$ -maximal subsets other than the ones described in Theorem 1;
- (ii) If  $d = n - 1$  then  $\mathcal{M}^n$  is the collection of all  $d$ -maximal subsets.

## 2 Proofs of the Theorems

Let  $A \subseteq B^n$ . Denote

$$\begin{aligned} G_t(A) &= \{\alpha \in B^n \setminus A \mid \rho(\alpha, A) \leq t\}, \\ \overline{A} &= \{\alpha \in B^n \mid \alpha = \overline{\beta}, \beta \in A\}. \end{aligned}$$

**Lemma 1**  $D(A) \leq d$  iff  $A \cap (\overline{A} \cup G_{n-d-1}(\overline{A})) = \emptyset$ .

*Proof.*

Assume  $D(A) \leq d$  and  $\alpha \in A$ . Then  $S_{n-d-1}^n(\overline{\alpha}) \cap A = \emptyset$ , because for  $\beta \in S_{n-d-1}^n(\overline{\alpha}) \cap A$  one has  $d + 1 \leq \rho(\alpha, \beta) \leq D(A)$ . Note that

$$\overline{A} \cup G_{n-d-1}(\overline{A}) = \bigcup_{\alpha \in \overline{A}} S_{n-d-1}^n(\alpha),$$

which implies the necessity of the above condition.

To prove its sufficiency, assume  $A \cap (\overline{A} \cup G_{n-d-1}(\overline{A})) = \emptyset$ . If  $D(A) \geq d + 1$ , then there exist vectors  $\alpha, \beta \in A$  such that  $\rho(\alpha, \beta) \geq d + 1$ . Therefore,  $\beta \in S_{n-d-1}^n(\alpha)$  and  $A \cap (\overline{A} \cup G_{n-d-1}(\overline{A})) \neq \emptyset$ . This contradiction completes the proof.  $\square$

**Corollary 1** If  $A \subseteq B^n$  is a  $d$ -maximal subset, then  $2|A| + |G_{n-d-1}(A)| \leq 2^n$ .  $\square$

Denote by  $I_t^n(m)$  the problem that consists in finding an  $m$ -element subset  $A_0 \subseteq B^n$  such that  $|G_t(A_0)| \leq |G_t(A)|$  for any  $A \subseteq B^n$ ,  $|A| = m$ . Denote  $g(m, n, t) = |G_t(A_0)|$ . We show that the function  $f(m, n, t) = 2m + g(m, n, t)$  is non-decreasing in  $m$  for fixed  $n, t$ . Indeed, the inequality  $f(m, n, t) < f(m + 1, n, t)$  is equivalent to  $g(m, n, t) \leq 1 + g(m + 1, n, t)$ . On the other hand, if  $C \subseteq B^n$  is a solution to  $I_t^n(m + 1)$  and  $\alpha \in C$ , then  $g(m, n, t) \leq |G_t(C \setminus \alpha)| \leq g(m + 1, n, t) + 1$ .

*Proof of Theorem 1.*

Assume  $d = 2t$ . It is known from [3] that  $S_t^n(\alpha)$  is a solution to  $I_s^n(m_0)$  for any  $s = 1, \dots, n$ , where

$$m_0 = |S_t^n(\alpha)| = \sum_{i=0}^t \binom{n}{i}.$$

One has

$$2m_0 + g(m_0, n, n-d-1) = \sum_{i=0}^t \binom{n}{i} + \sum_{i=0}^{n-t-1} \binom{n}{i} = 2^n.$$

Let  $A$  be a  $d$ -maximal subset and  $|A| = m$ . By Corollary 1,  $2m + |G_{n-d-1}(A)| \leq 2^n$ , hence  $2m + g(m, n, n-d-1) \leq 2^n$ . Using the monotonicity of  $g$  and the equality  $D(S_t^n(\alpha)) = d$  one gets  $m = m_0$ .

Now, assume  $d = 2t + 1$ . It is known from [3] that  $S_t^n(\alpha) \cup S_t^n(\beta)$  with  $\rho(\alpha, \beta) = 1$  is a solution to  $I_s^n(m_0)$  for any  $s = 1, \dots, n$ , where

$$m_0 = |S_t^n(\alpha) \cup S_t^n(\beta)| = \sum_{i=0}^t \binom{n}{i} + \binom{n-1}{t}.$$

Taking into account the equality  $2m_0 + g(m_0, n, n-d-1) = 2^n$  and using the arguments above, one gets that if  $A$  is a  $d$ -maximal subset then  $|A| = m_0$ , which completes the proof of the Theorem.  $\square$

**Corollary 2** *If  $A$  is a  $d$ -maximal subset, then  $A$  is a solution to  $I_{n-d-1}^n(|A|)$ .*  $\square$

**Lemma 2** *Let  $d < n - 1$  and  $A$  is a  $d$ -maximal subset. Then  $A$  is a solution to  $I_1^n(|A|)$ .*

*Proof.*

It follows from the proof of Theorem 1, that  $|A| + |\overline{A} \cup G_{n-d-1}(\overline{A})| = 2^n$ . Since  $A \cap (\overline{A} \cup G_{n-d-1}(\overline{A})) = \emptyset$  then

$$A = B^n \setminus (\overline{A} \cup G_{n-d-1}(\overline{A})).$$

Therefore,  $|G_t(A)| = |G_{n-d-1}(\overline{A})| - |G_{n-d-2}(\overline{A})|$ .

Let  $d = 2t$ . One has

$$\begin{aligned} |A| &= \sum_{i=0}^t \binom{n}{i}, \\ |G_{n-d-1}(\overline{A})| &= 2^n - 2|A| = \sum_{i=1}^{n-d-1} \binom{n}{t+i}, \\ |G_{n-d-2}(\overline{A})| &\geq \sum_{i=1}^{n-d-2} \binom{n}{t+i}. \end{aligned}$$

This implies  $|G_1(A)| \leq \binom{n}{t+1}$ , i.e.  $A$  is a solution to  $I_1^n(|A|)$ .

The proof of Lemma in the case  $d = 2t + 1 < n - 1$  is similar.  $\square$

*Proof of Theorem 2.*

Let  $A \subseteq B^n$  be a  $d$ -maximal subset. Denote by  $I^n(m)$  the problem that consists in finding an  $m$ -element subset  $A_0 \subseteq B^n$ , such that  $|\Gamma(A_0)| \leq |\Gamma(A)|$  for any  $A \subseteq B^n$ ,  $|A| = m$ , where

$$\Gamma(A) = \{\alpha \in A \mid S_1^n(\alpha) \notin A\}.$$

Since  $G_1(A) = \Gamma(B^n \setminus A)$ , the set of solutions to  $I^n(m)$  can be obtained from the one for  $I_1^n(2^n - m)$  and vice versa. In particular, denoting  $C = B^n \setminus A$ , we get by Lemma 2 that  $C$  is a solution to  $I^n(2^n - |A|)$ .

Assume  $d = 2t < n - 1$ . Denote

$$m_1 = |A| = \sum_{i=0}^t \binom{n}{i}, \quad m_2 = |C| = \sum_{i=0}^{n-t-1} \binom{n}{i}.$$

It is proved in [3] (Theorem 6.1), that  $S_{n-t-1}^n(\alpha)$  is the only solution to  $I^n(m_2)$  up to the choice of  $\alpha \in B^n$ . Hence,  $S_t^n(\bar{\alpha})$  is the only solution to  $I_1^n(m_1)$ . Therefore,  $A = S_t^n(\bar{\alpha})$ .

Assume  $d = 2t + 1 < n - 1$ . Denote

$$m_3 = |A| = \sum_{i=0}^t \binom{n}{i} + \binom{n-1}{t}, \quad m_4 = |C| = \sum_{i=0}^{n-t-2} \binom{n}{i} + \binom{n-1}{t+1}.$$

It is proved in [3] (Lemma 6.1) that there exists a coordinate  $x_j$ ,  $1 \leq j \leq n$ , such that  $C^0(j) = S_{n-t-2}^{n-1}(\alpha')$  and  $C^1(j) = S_{n-t-2}^{n-1}(\beta')$ , where  $C^0(j)$  and  $C^1(j)$  are the intersections of the set  $C$  with the hyperplanes  $x_j = 0$  and  $x_j = 1$  respectively. Here  $\alpha'$  and  $\beta'$  are projections of resp.  $\alpha$  and  $\beta$  to the correspondent hyperplanes. By Theorem 6.2 of [3], we have  $\rho(\alpha, \beta) \leq 2$ . It is not difficult to show that if  $\rho(\alpha, \beta) = 2$ , then  $D(A) > d$ . This implies,  $\rho(\alpha, \beta) = 1$ , i.e the vectors  $\alpha, \beta$  differ in the  $j$ -th entry only. Furthermore,

$$C = S_{n-t-2}^{n-1}(\alpha) \cup S_{n-t-2}^{n-1}(\beta) = S_{n-t-2}^n(\alpha) \cup S_{n-t-2}^n(\beta).$$

Therefore  $A = S_t^n(\bar{\alpha}) \cup S_t^n(\bar{\beta})$ .

The proof in the case  $d = n - 1$  follows from Lemma 1. □

### 3 Applications of the Theorems

**1.** Consider the problem of constructing an  $m$ -element subset of  $B^n$  with minimal possible diameter. Represent the integer  $m$  in the form

$$m = \sum_{i=0}^k \binom{n}{i} + \delta, \quad 0 < \delta \leq \binom{n}{k+1}.$$

Obviously,  $D(A) = n$  for any  $A \subseteq B^n$ ,  $|A| \geq 2^{n-1}$ . If  $m \leq 2^{n-1}$  and  $\binom{n-1}{k} < \delta \leq \binom{n}{k+1}$ , Theorem 1 implies  $D(A) \geq 2k + 2$  for any  $A \subseteq B^n$ ,  $|A| = m$ . However, such a subset  $A$  may be obtained by deleting arbitrary  $\sum_{i=0}^{k+1} \binom{n}{i} - m$  vectors from a ball of radius  $k + 1$ . The argument in the case  $m \leq 2^{n-1}$  and  $0 < \delta < \binom{n-1}{k}$  is similar.

**2.** Consider arbitrary  $n$ -element set and let  $\mathcal{F}_l$  be a family of it's subsets such that  $|F_i \cap F_j| \geq l > 1$  for any  $F_i, F_j \in \mathcal{F}_l$ ,  $i \neq j$ . We call such a family, consisting of maximal possible number of subsets, the maximal one. A complete specification of the maximal

families with  $l > 1$  can be found in [4]. A solution to this problem can also be deduced from Theorem 2.

For this note, that there is a natural correspondence between the vectors of the  $n$ -cube  $B^n$  and the subsets of  $n$ -element set and the Hamming distance between two vertices of the  $n$ -cube is equal to the size of the symmetric difference of the corresponding subsets.

If  $F_i, F_j \in \mathcal{F}_l$ , then  $|F_i \Delta F_j| = |F_i \cup F_j| - |F_i \cap F_j| \leq n - l$ . Therefore, if  $n - l = 2t$  then  $D(\mathcal{F}_l) \leq 2t$  and, by Theorem 1,

$$|\mathcal{F}_l| \leq \sum_{i=0}^t \binom{n}{i} = m_0.$$

On the other hand, the ball  $S_t^n(\mathbf{1})$ , where  $\mathbf{1} = (1, \dots, 1) \in B^n$ , corresponds to some family  $\mathcal{F}_l$ . Hence, the size of any maximal family equals  $m_0$ . Similarly, if  $n - l = 2t + 1$  is odd, then one can easily show that the diameter of any maximal family  $\mathcal{F}_l$  is equal to  $n - l$ . Therefore, any maximal family corresponds to some  $(n - l)$ -maximal subset, i.e. for some  $\alpha, \beta, \gamma \in B^n$   $\mathcal{F}_l = S_t^n(\alpha)$  for  $n - l$  even and  $\mathcal{F}_l = S_t^n(\beta) \cup S_t^n(\gamma)$  with  $\rho(\beta, \gamma) = 1$  for  $n - l$  odd. Furthermore, it is easy to verify that if  $\alpha \neq \mathbf{1}$  and  $\beta \neq \mathbf{1}$  (or  $\gamma \neq \mathbf{1}$ ), then there always exist some  $F_i, F_j \in \mathcal{F}_l$  such that  $|F_i \cap F_j| < l$ . Therefore if  $n - l$  is even, then there exists the only maximal family and there are exactly  $n$  maximal families for  $n - l$  is odd.

Note that for  $l = 1$  the above statements are not valid. It is not difficult to show [5] that in this case for the number  $M$  of maximal families one has

$$\log_2 M \geq 0.5 \binom{n}{\lfloor n/2 \rfloor}, \quad \text{as } n \rightarrow \infty.$$

Following a conjecture in [5],  $\log_2 M \sim 0.5 \binom{n}{\lfloor n/2 \rfloor}$  as  $n \rightarrow \infty$ .

## References

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